

From Dataflow to Circuits

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Outline

- * Circuit implementation style
 - MOUSETRAP
- * Design example
 - Greatest-Common Divisor (GCD)

Implementation of asynchronous pipelines

MOUSETRAP Circuit Style

MOUSETRAP Pipelines

[Singh/Nowick, ICCD 2001 & TVLSI 2007]

Simple asynchronous implementation style, uses...

- *transparent D-latches + standard combinational function logic*
- *simple control:* 1 gate/pipeline stage

Uses a “capture protocol”: Latches are ...

- normally transparent while waiting for data
- become opaque after data arrives

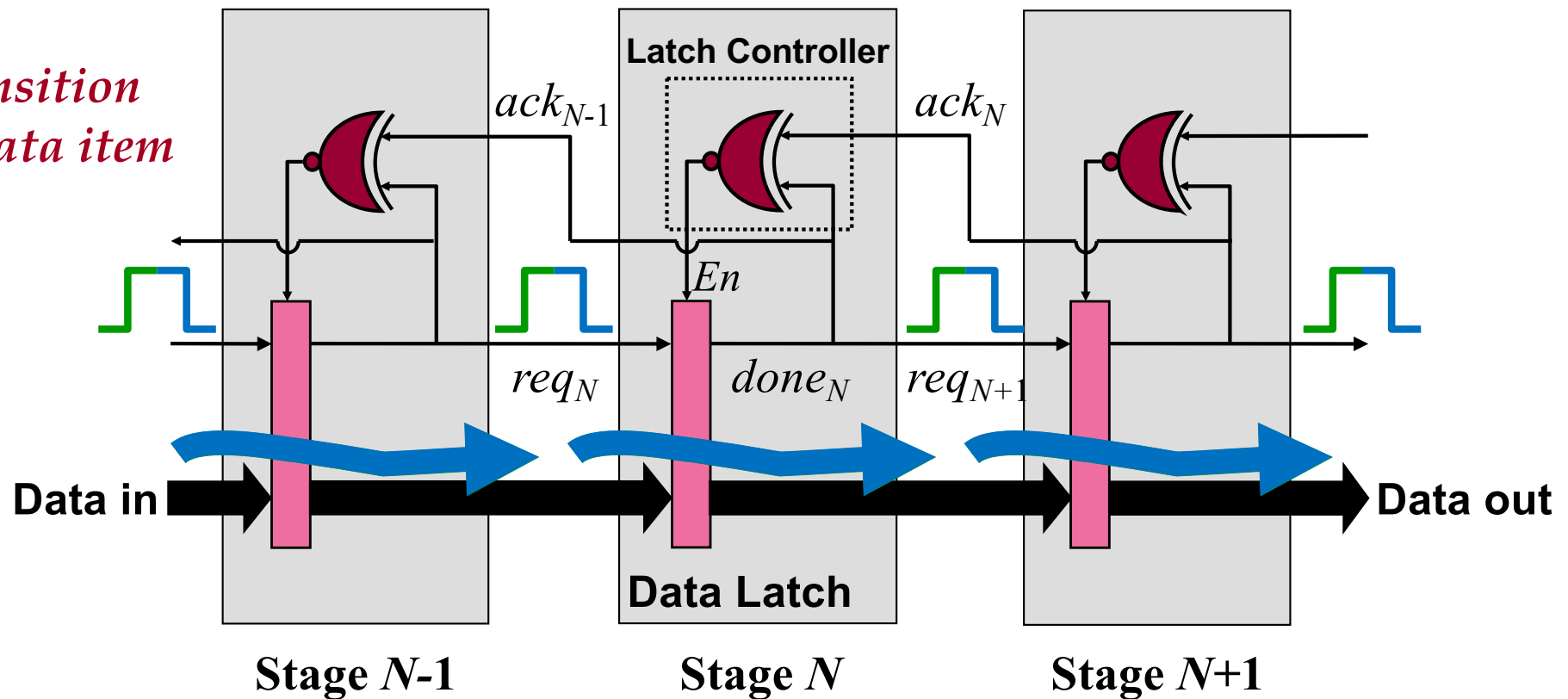
Control Signaling: *transition-signaling = 2-phase*

Goals:

- fast cycle time
- simple inter-stage communication
- standard cell implementation

MOUSETRAP: A Basic FIFO

Stages communicate using *transition-signaling*:

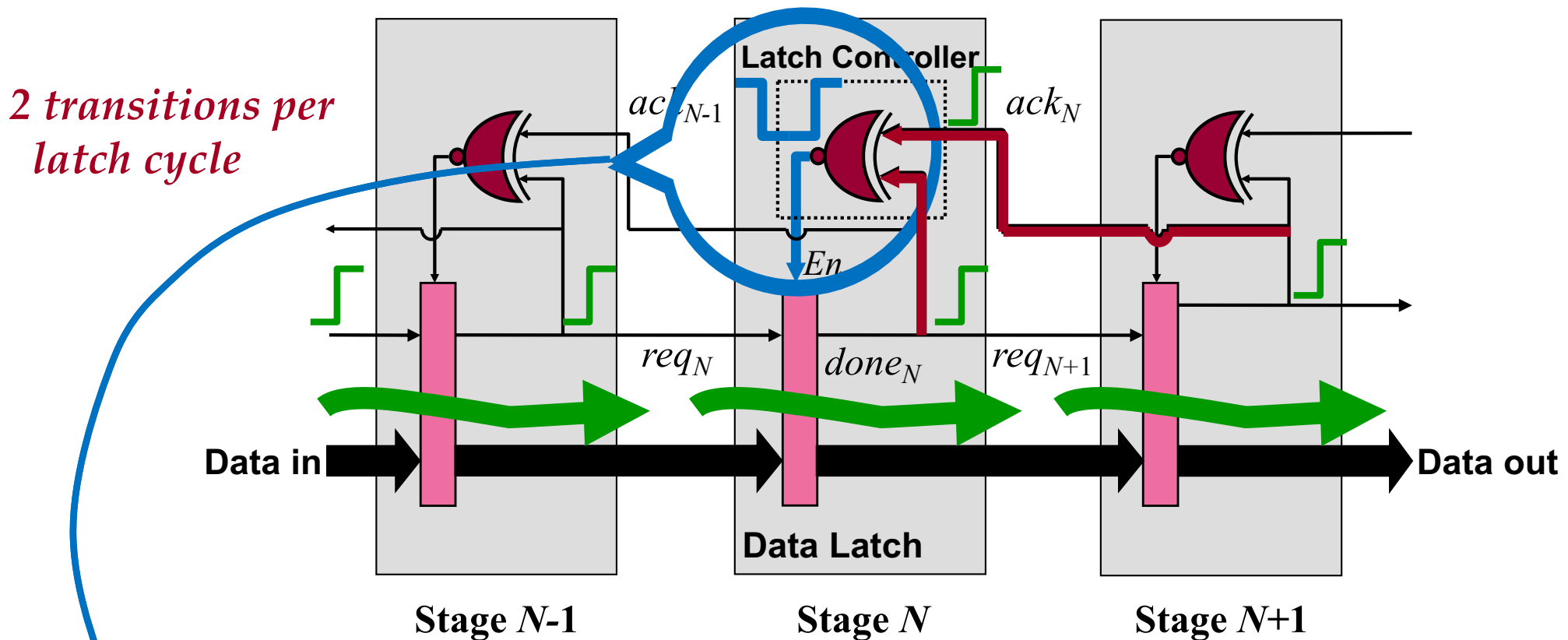


2nd data item flowing through the pipeline

MOUSETRAP: A Basic FIFO (contd.)

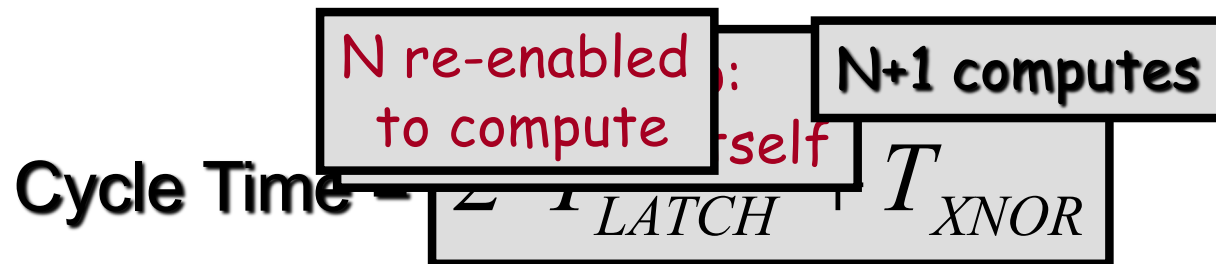
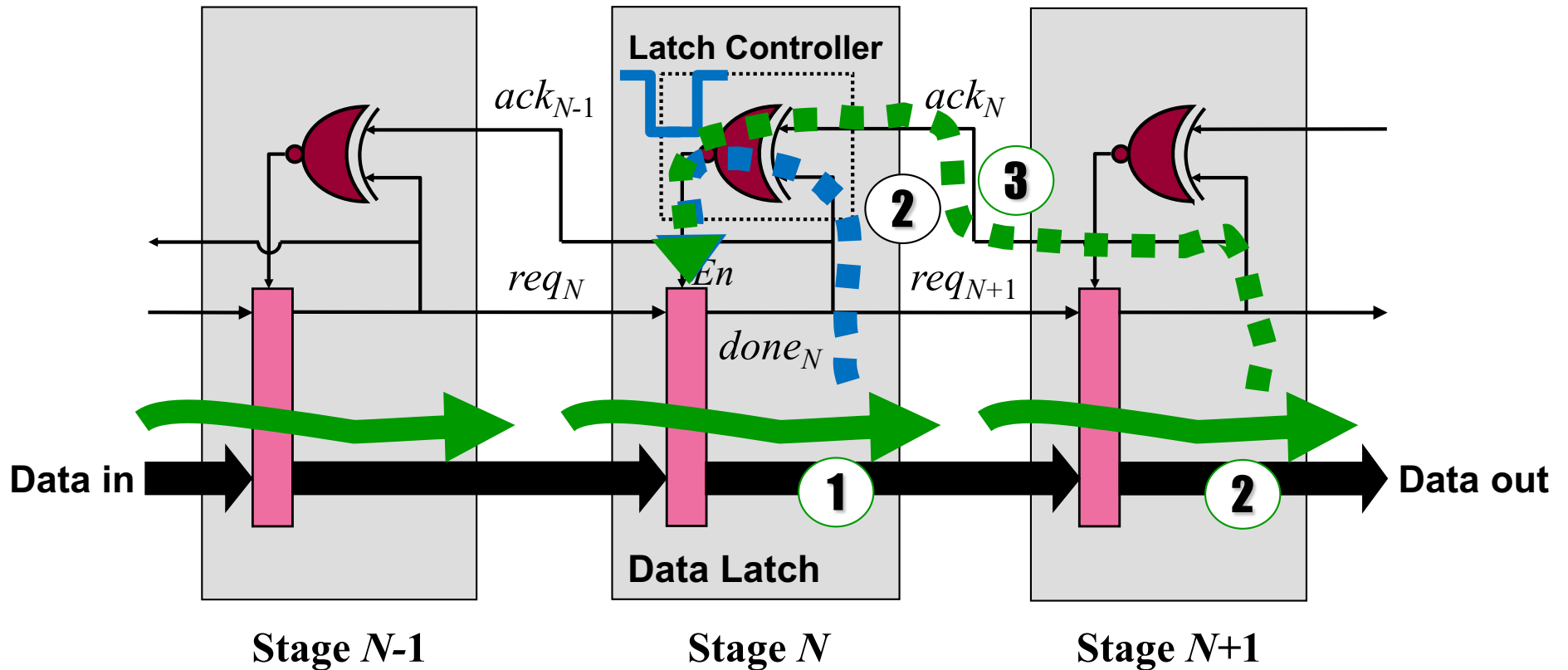
Latch controller (XNOR) acts as “*phase converter*”:

- 2 distinct transitions (up or down) → pulsed latch enable



Latch is disabled when next stage is “done”

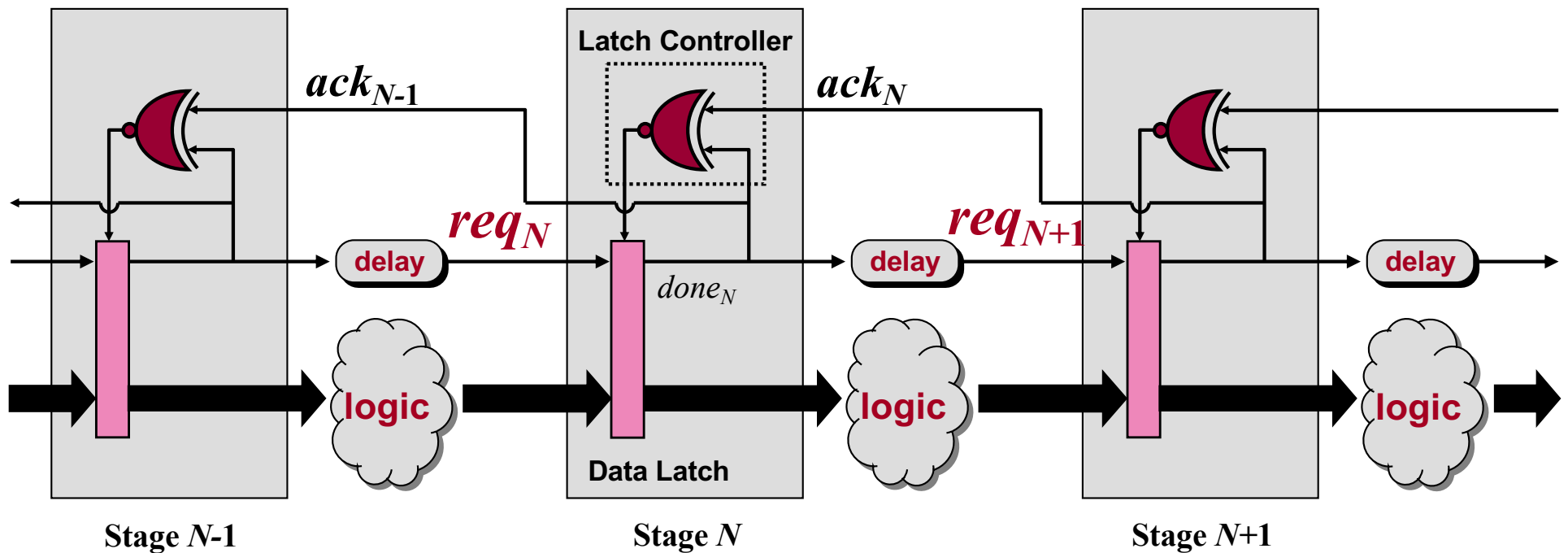
MOUSETRAP: FIFO Cycle Time



MOUSETRAP: Pipeline With Logic

Simple Extension to FIFO:

insert *logic block* + *matching delay* in each stage

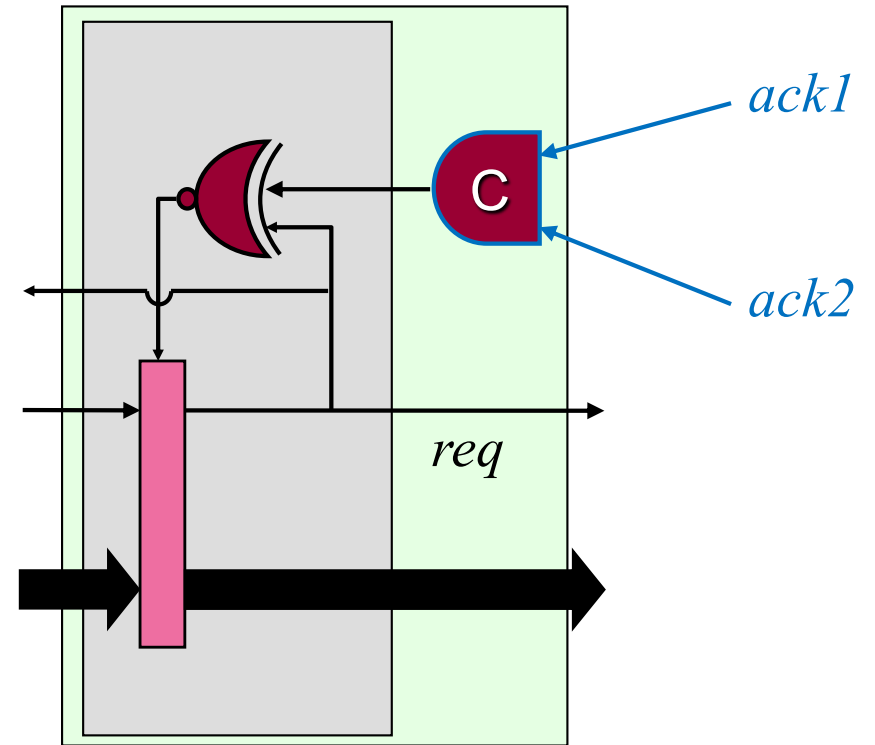


"Bundled Data" Requirement:

- each *req* must arrive *after* data inputs valid and stable

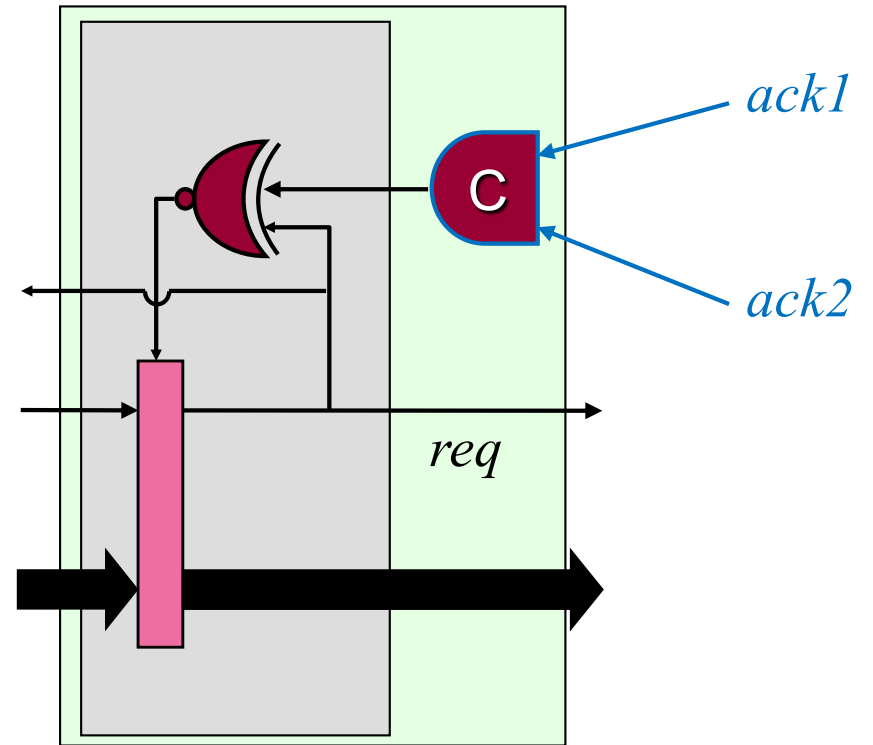
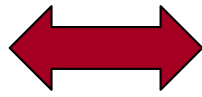
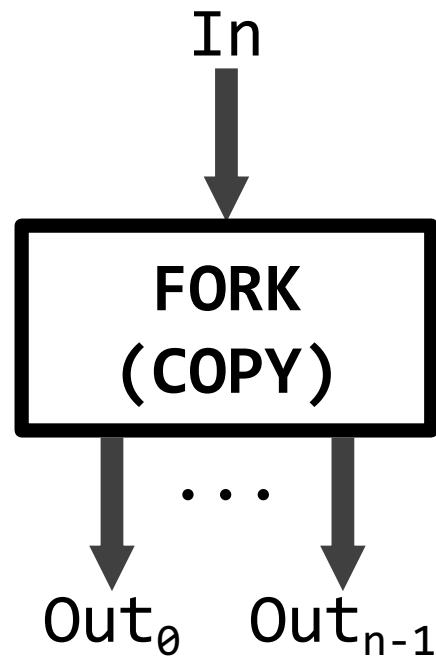
Fork

- * A fork stage has two (or more) successors
 - same data and req sent to all
 - wait for ack from all



Fork Stage

Fork



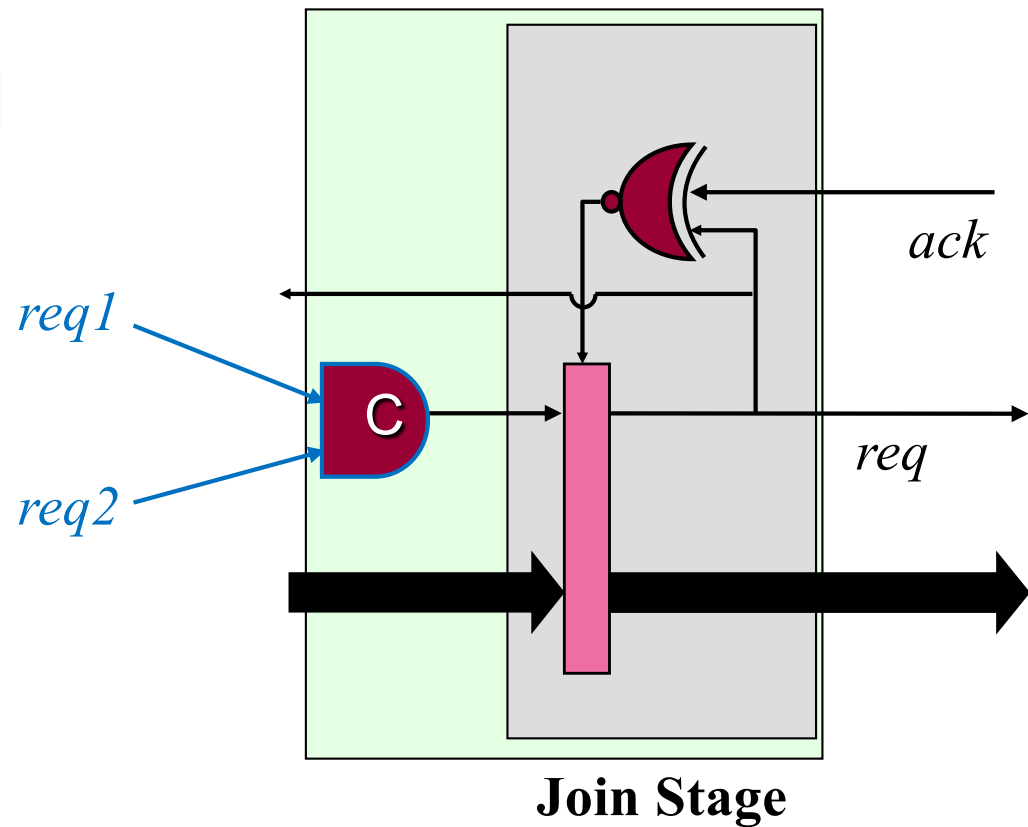
Fork Stage

*[In?x; Out₀!x, ..., Out_{n-1}!x]

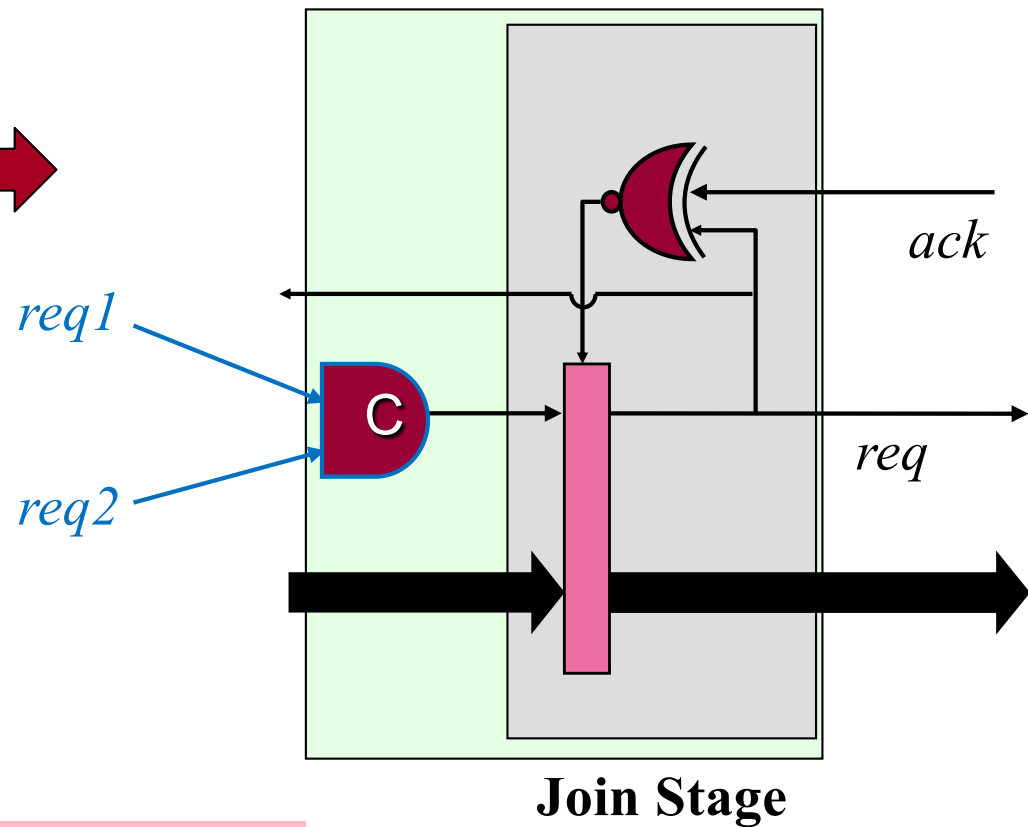
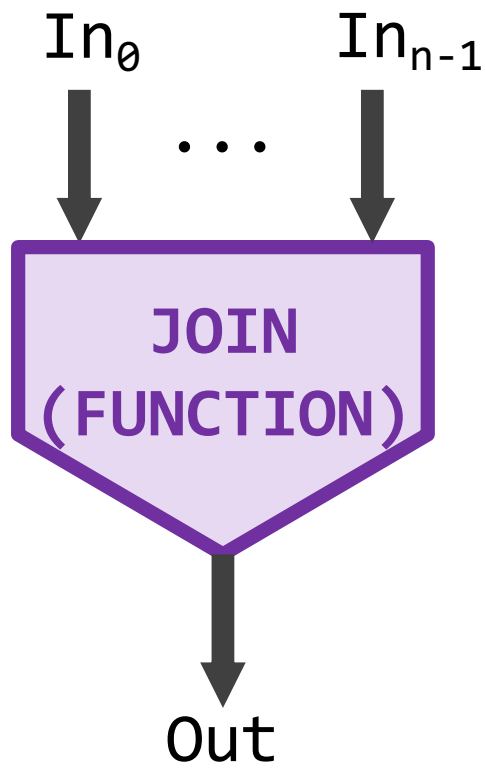
Join

* A join stage has two (or more) predecessors

- wait for data from all
- same ack sent to all



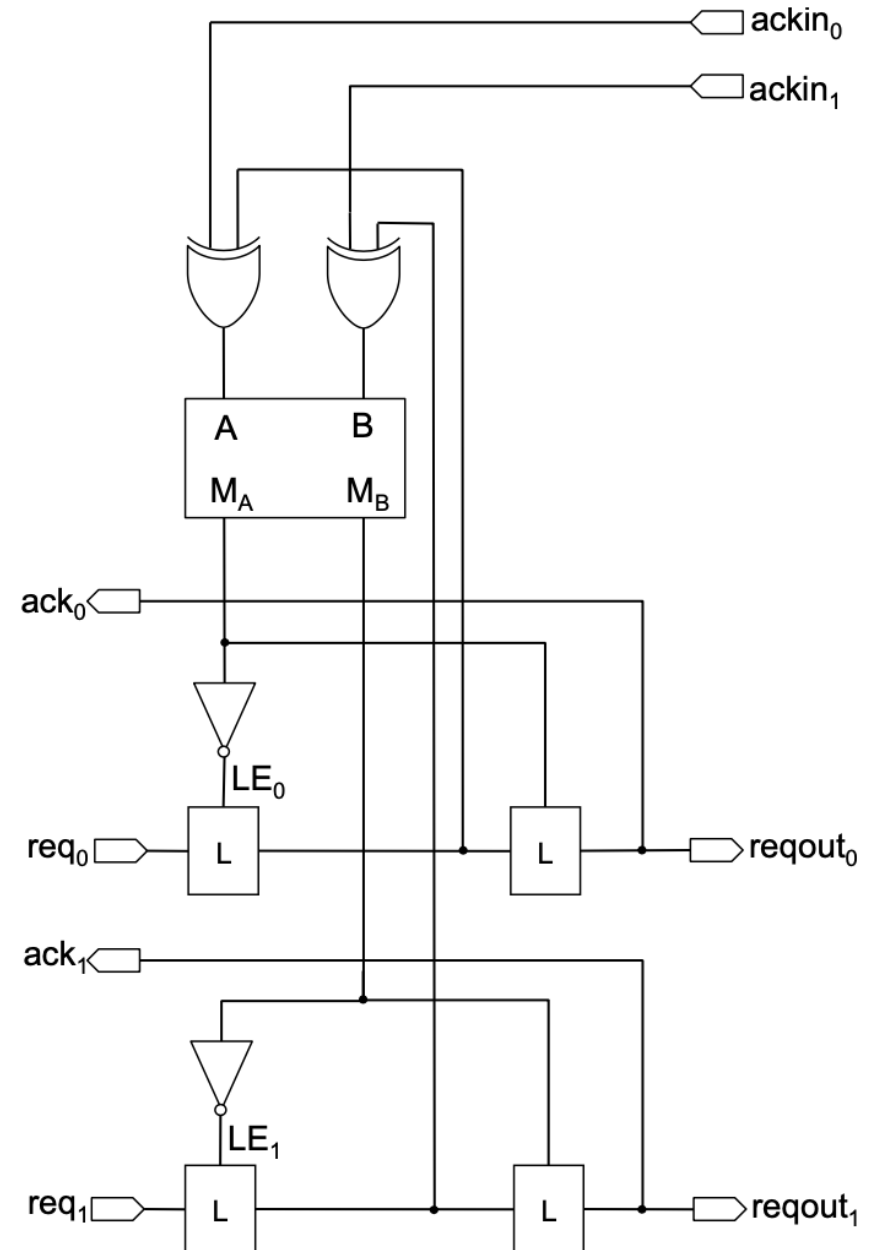
Join



```
*[ In0?arg0, In1?arg1, ... , Inn-1?argn-1;  
  Out!func(arg0,arg1,...,argn-1)  
]
```

Arbitration Stage

- * Two input channels, two output channels
- * Only one input read
 - whichever arrives first
 - ... goes out on the corresponding output
- * Seitz' Mutex is the core
 - [from Brunvand's thesis]
 - surrounding logic adapts it to transition signals



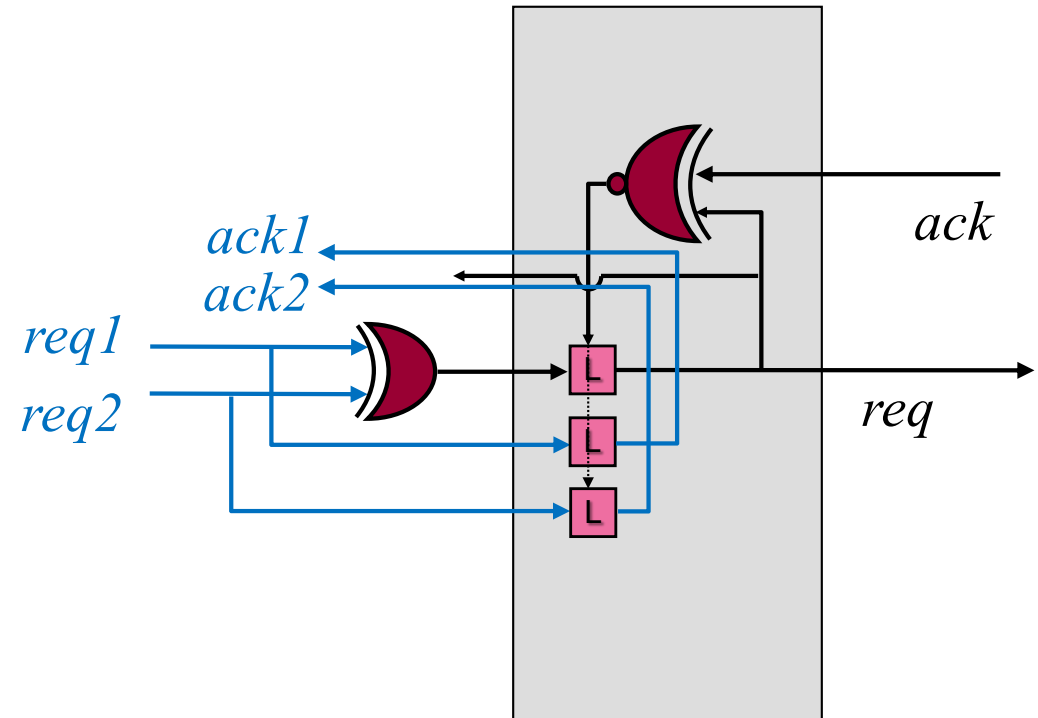
Merge without (or after) Arbitration

* A merge stage has two (or more) predecessors

- data is taken from whichever input channel has a new request

* Assumption:

- no arbitration needed
- input channels are mutually exclusive

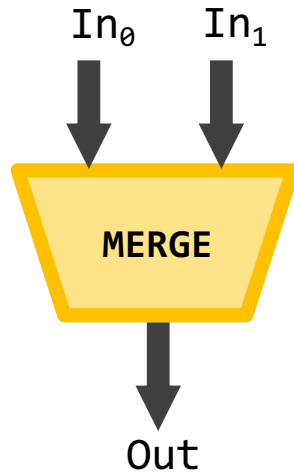


Merge w/o arbitration

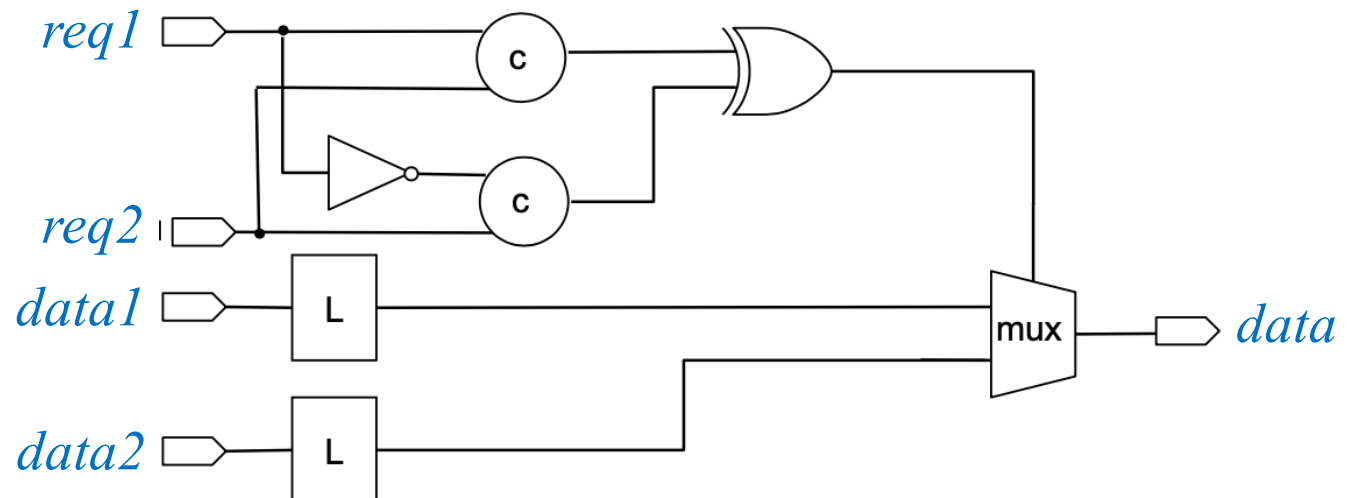
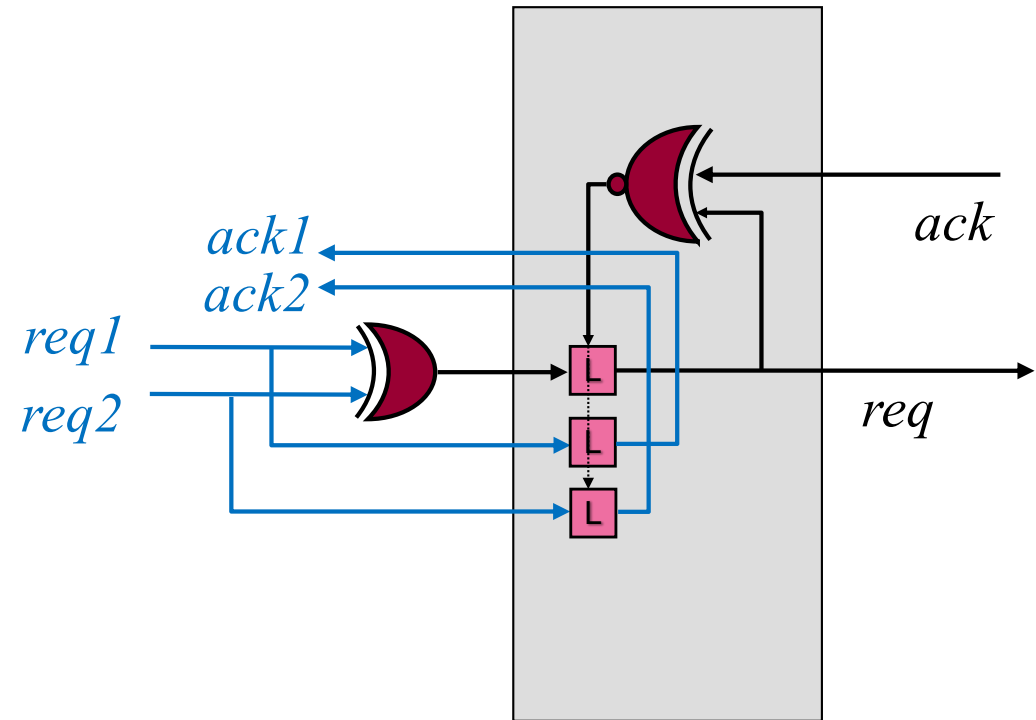
Merge without (or after) Arbitration

* Datapath

- mux controlled by change detectors on input channels



```
*[ [ #In_0 -> In_0?x
  [] #In_1 -> In_1?x
  ];
  Out!x
]
```



Conditional Select (or event mux)

* Two data inputs and one select

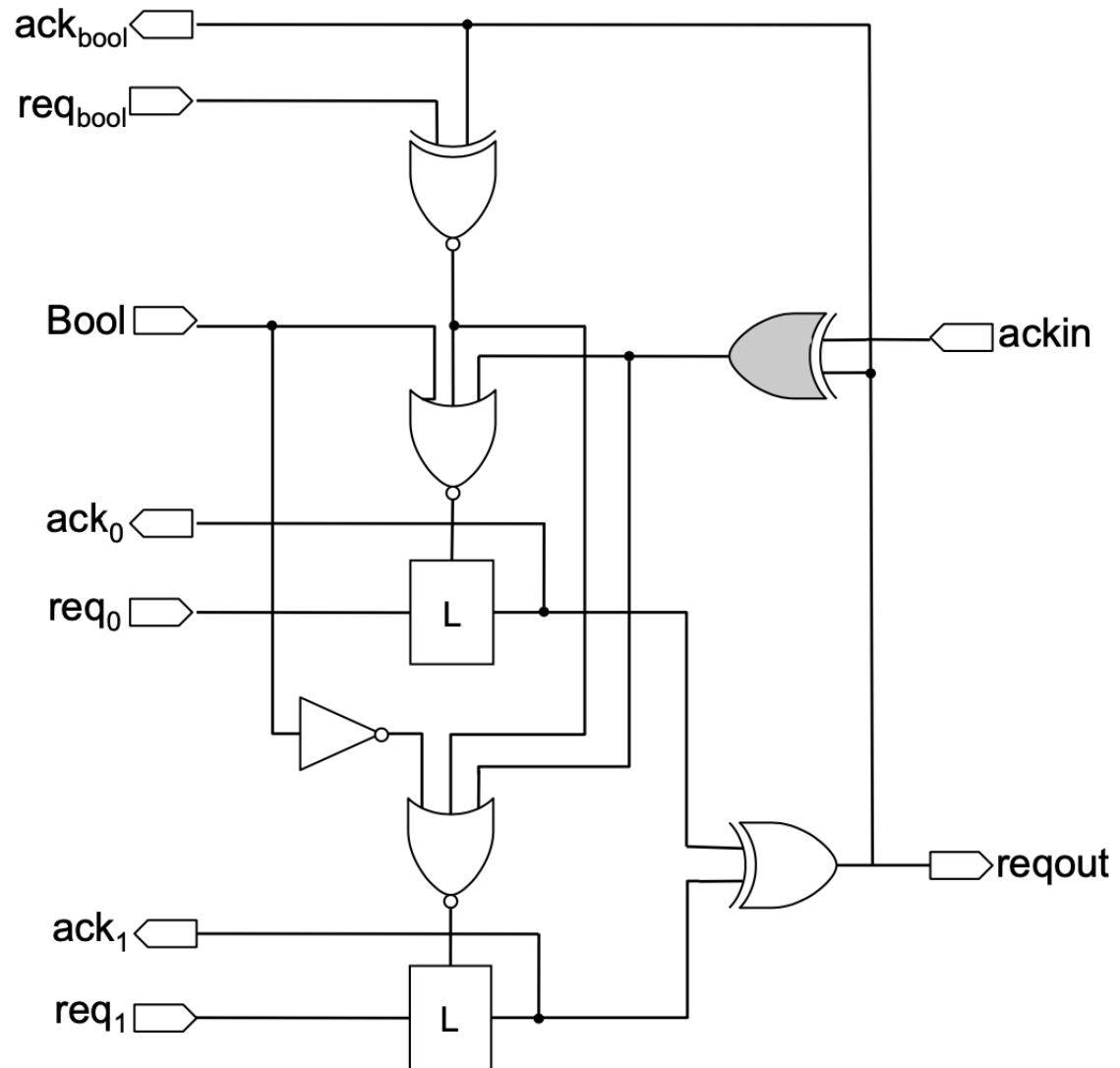
- first read select

- then:

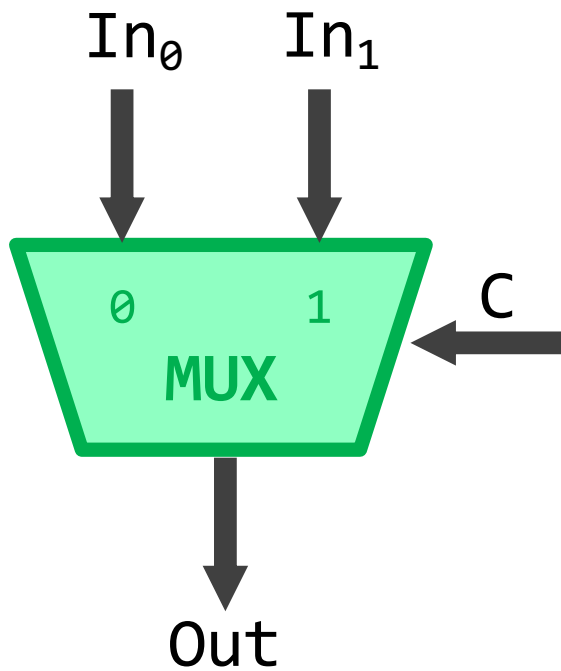
- based on select, read one of the inputs
- do not read the other input

- datapath

- mux + latch

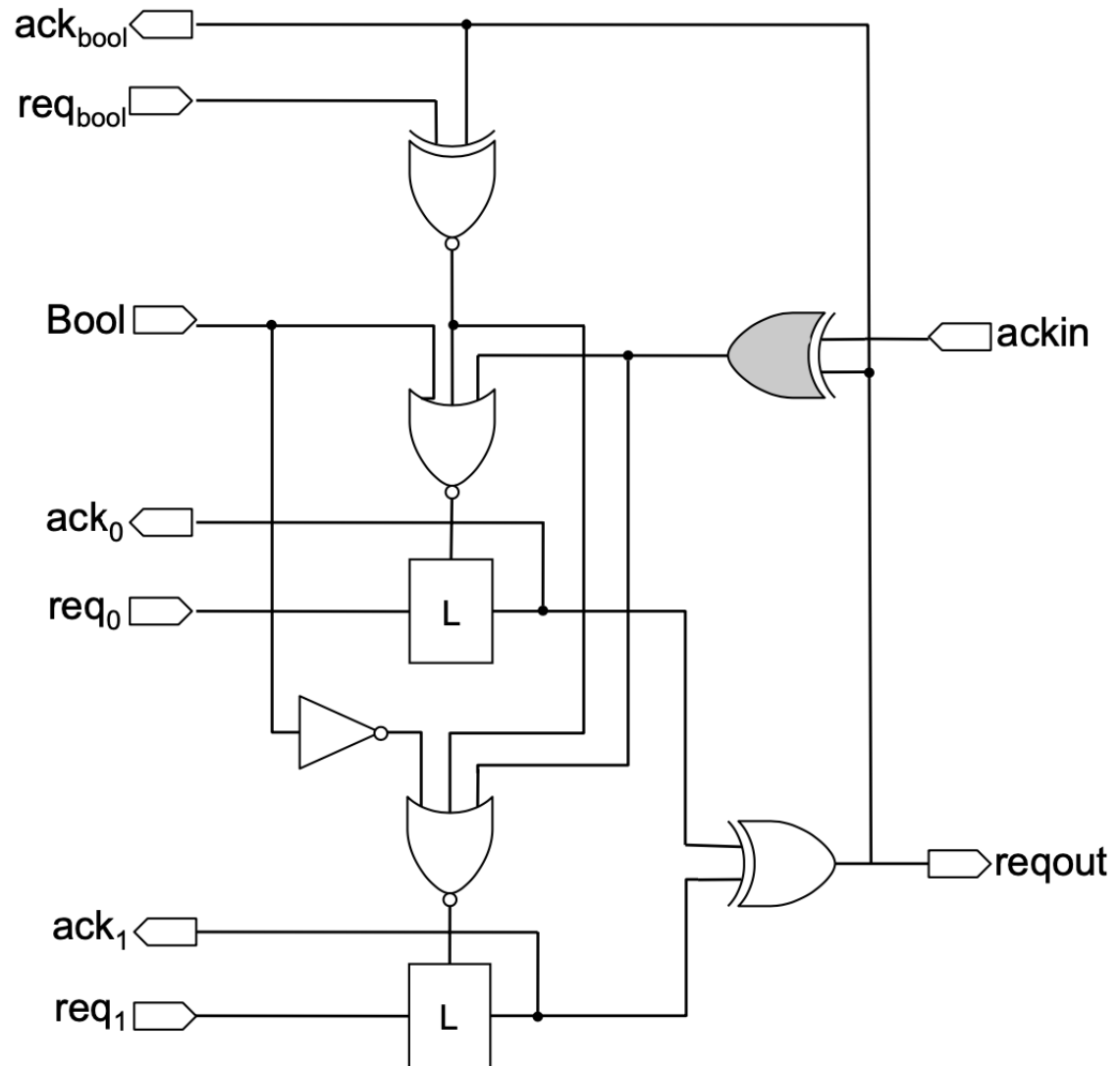


Conditional Select (or event mux)



```

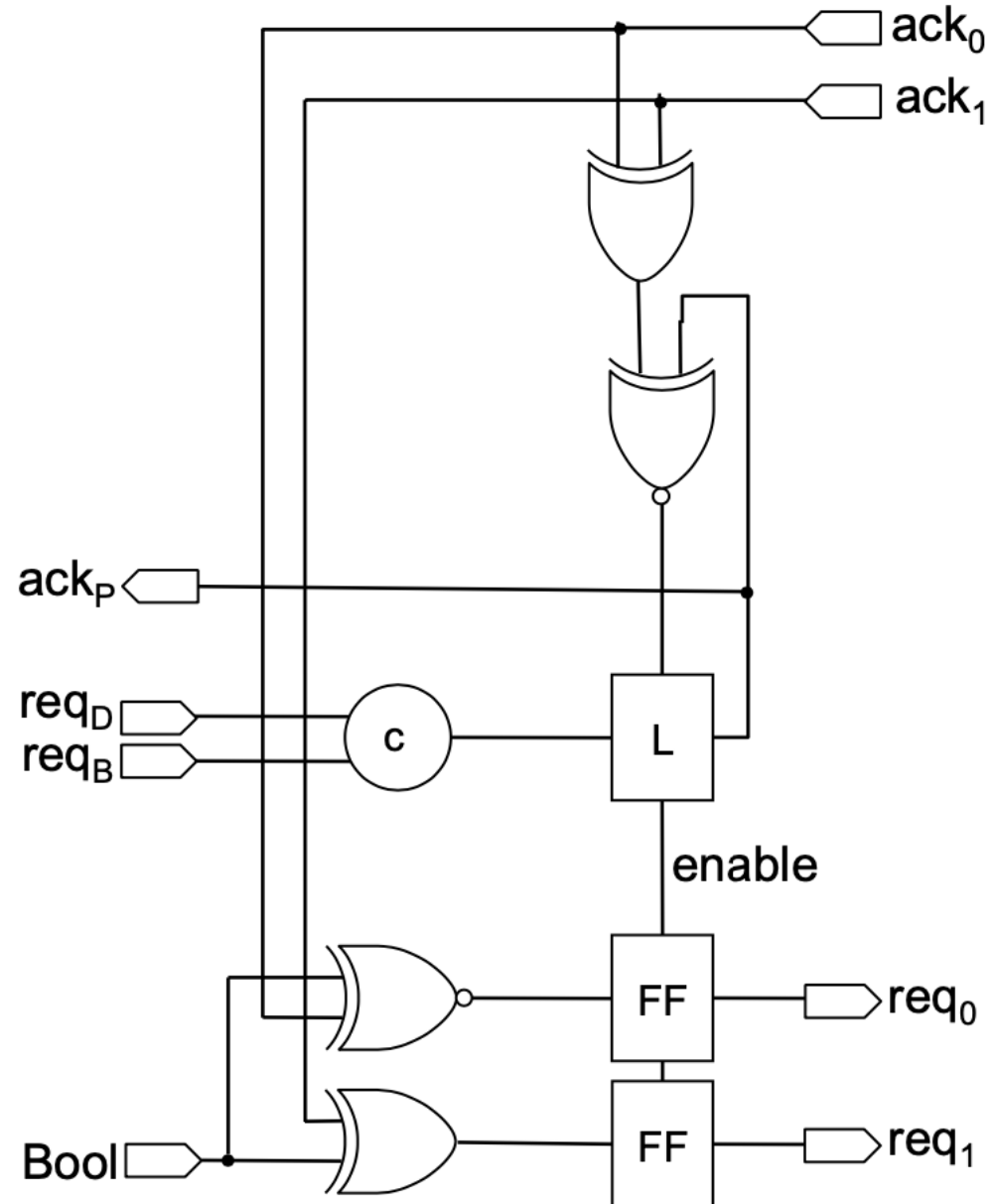
*[C?c;
 [ c=0 -> In_0?x
 [] c=1 -> In_1?x
 ];
 Out!x
 ]
    
```



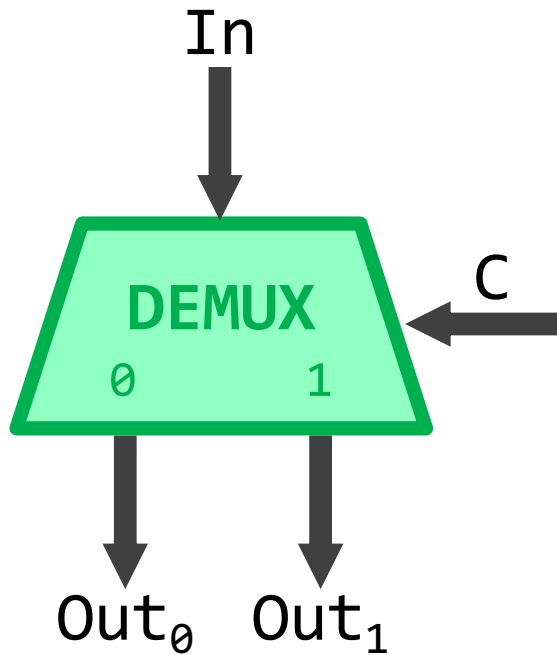
Conditional Split (or router)

* One data input and one select

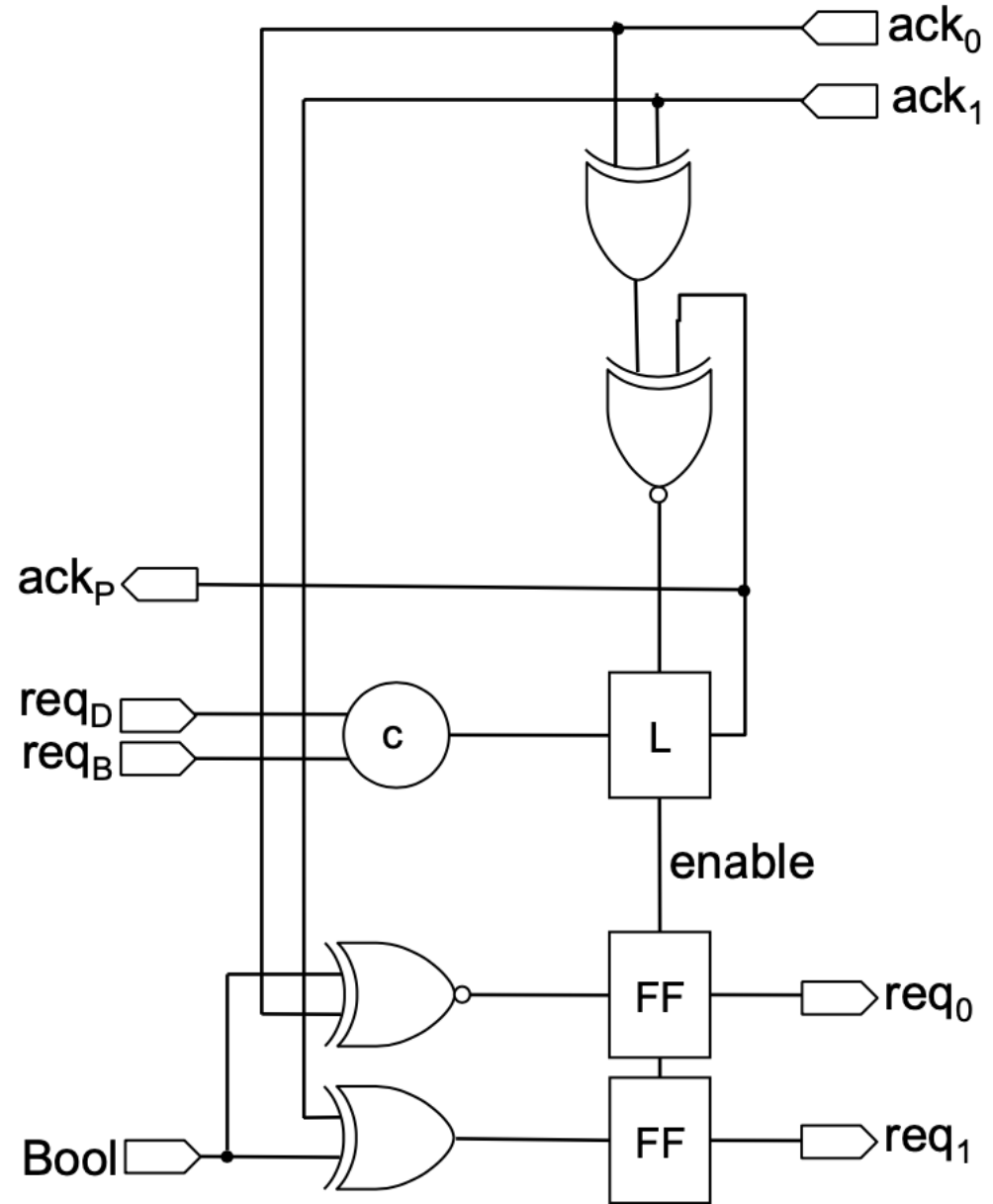
- first read data + select
- then:
 - based on select, send data along one output channel



Conditional Split (or router)



```
*[In?x, C?c;  
  [ c=0 -> Out0!x  
  [ ] c=1 -> Out1!x  
  ]  
]
```



Example: Greatest Common Divider

Euclid's GCD algorithm

```
gcd(a, b)
  while (b != 0)
    if(a>b)
      a = a - b
    else
      b = b - a
  return a
```

* Example

- gcd(42, 28)
- (14, 28)
- (14, 14)
- (14, 0)
- → 14

Euclid's GCD algorithm

`gcd(a, b)`

`while (b != 0)`

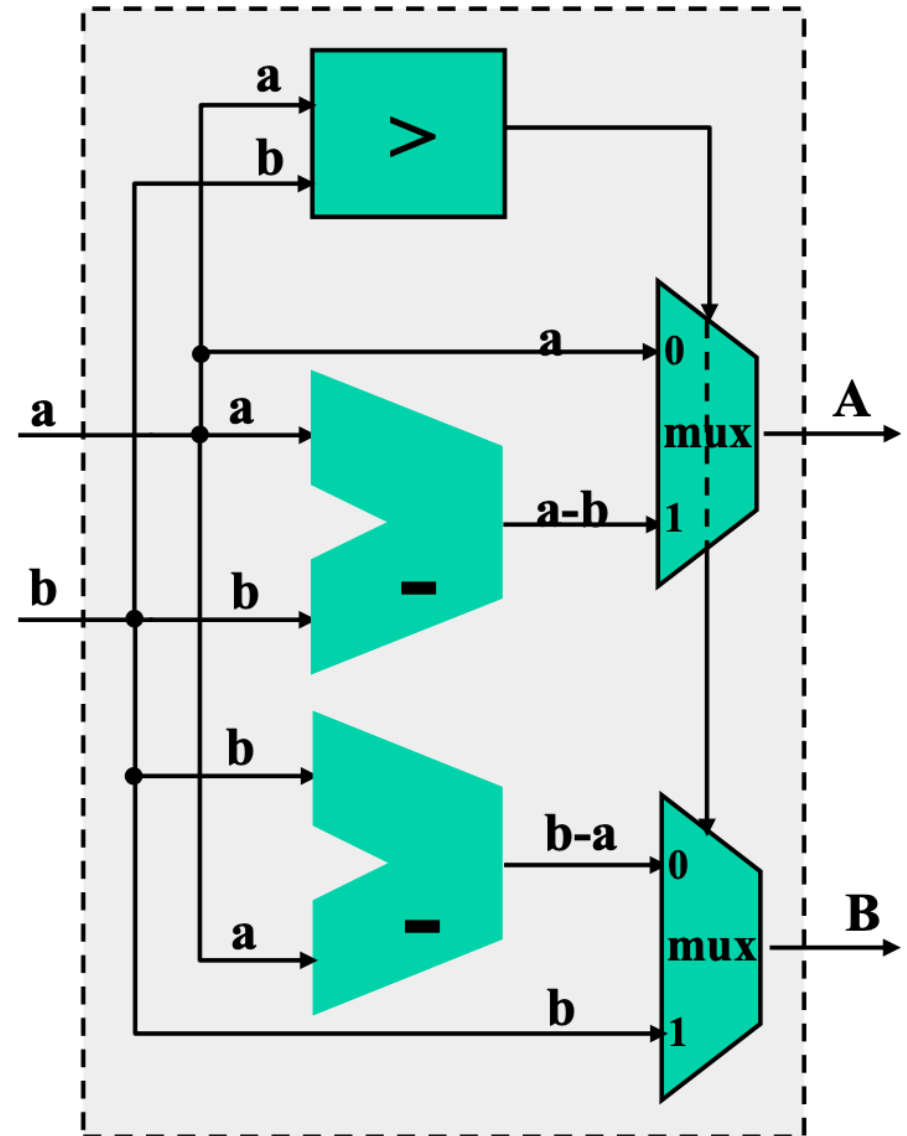
`if(a>b)`

`a = a - b`

`else`

`b = b - a`

`return a`



Better area version

gcd(a, b)

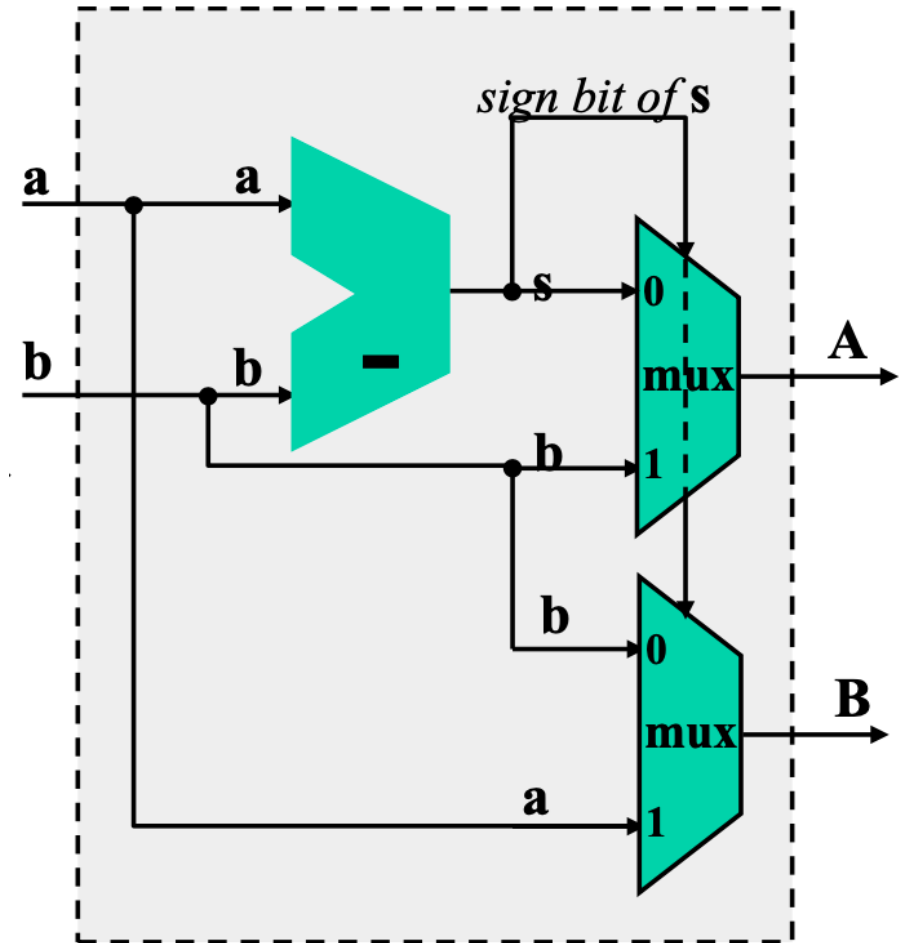
while (b != 0)

```
s = a-b
if(s<0)
    swap(a,b)
else
    a = s
```

return a

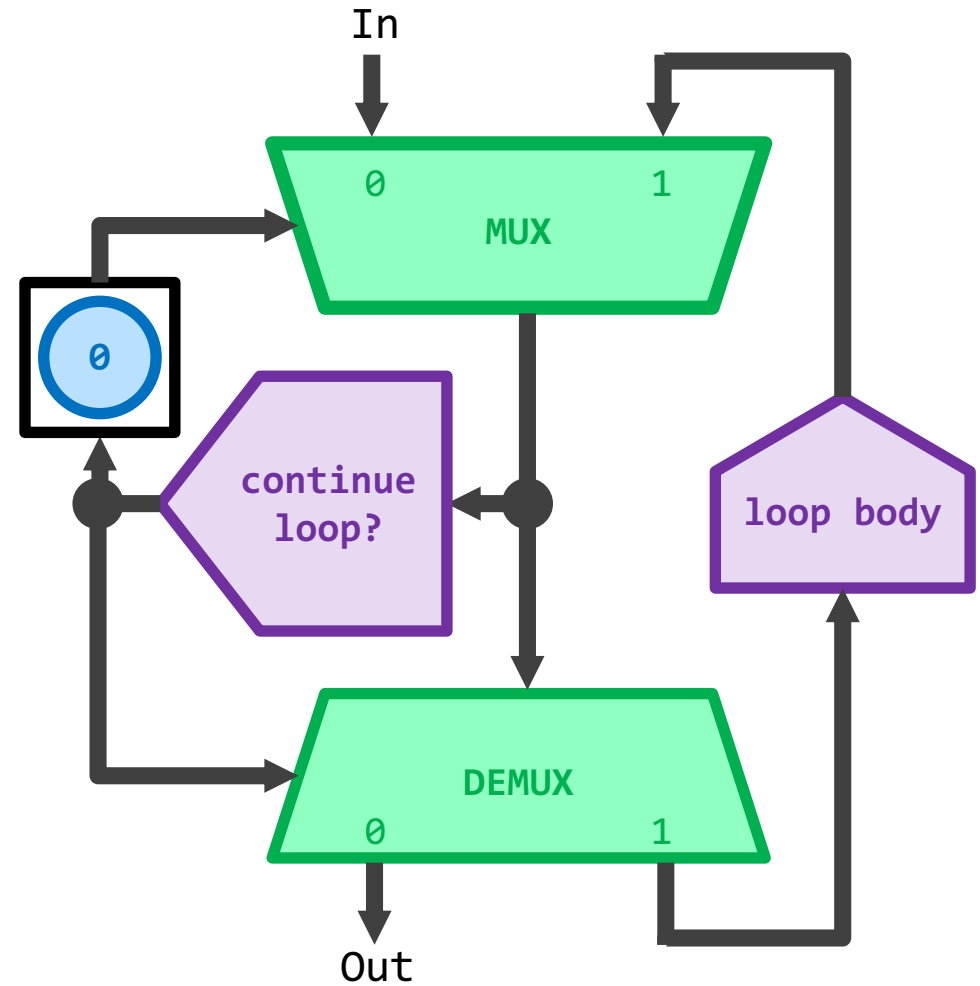
* Example

- gcd(42, 28)
- (14, 28)
- (28, 14)
- (14, 14)
- (0, 14)
- (14, 0)
- → 14



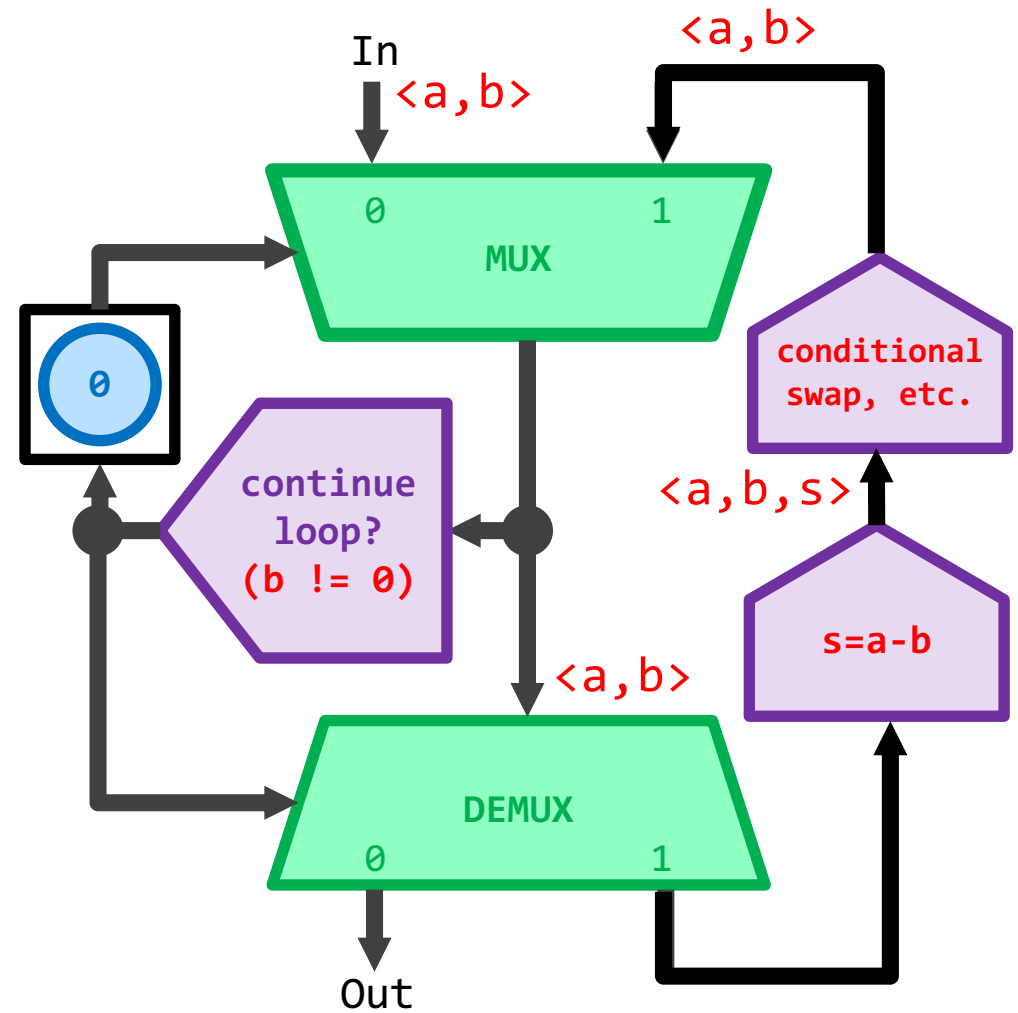
Dataflow: WHILE loop structure

```
while (b != 0)
  s = a-b
  if(s<0)
    swap(a,b)
  else
    a = s
return a
```



GCD implementation

```
while (b != 0)
  s = a-b
  if(s<0)
    swap(a,b)
  else
    a = s
return a
```



GCD implementation

```
while (b != 0)
```

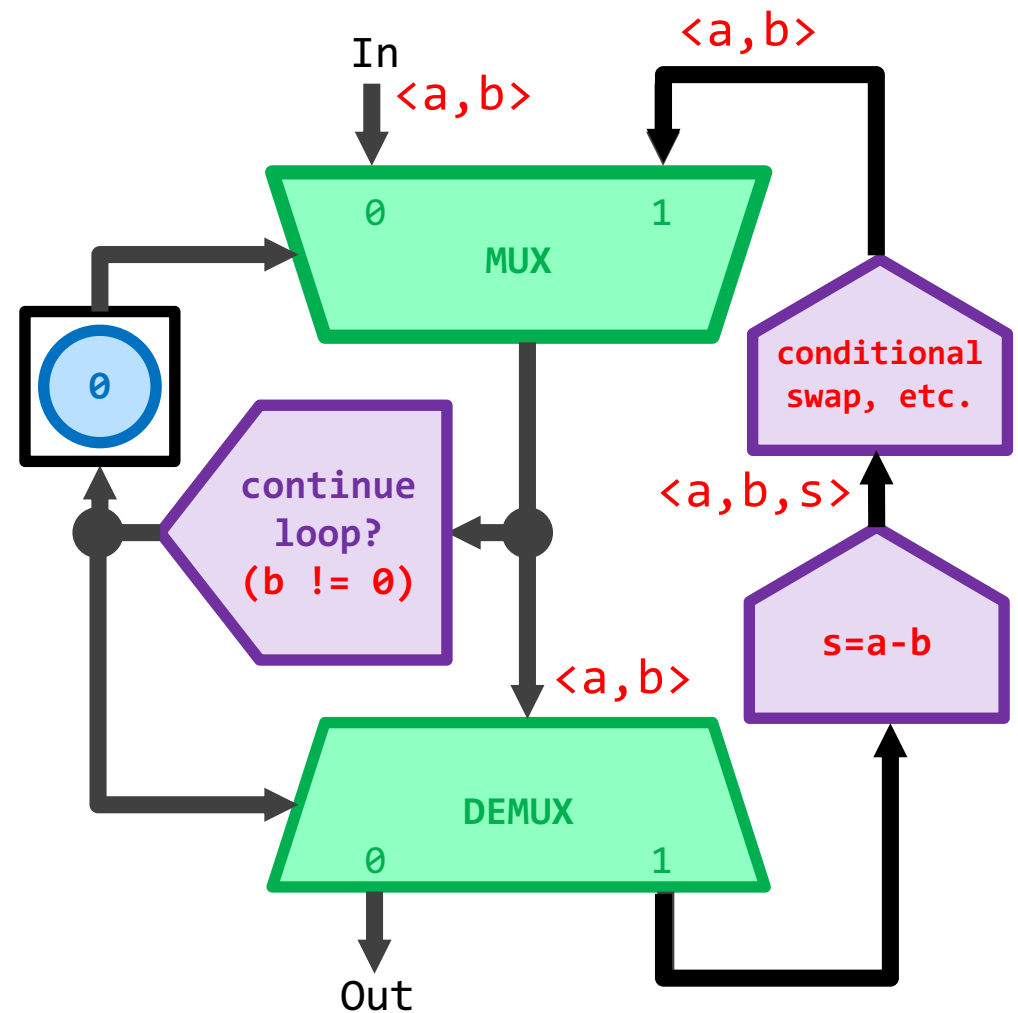
```
  s = a-b  
  if(s<0)  
    swap(a,b)  
  else  
    a = s
```

```
return a
```

Unroll 8 times!

Pipeline subtract
into 8 stages

Pipeline swap
into 4 stage



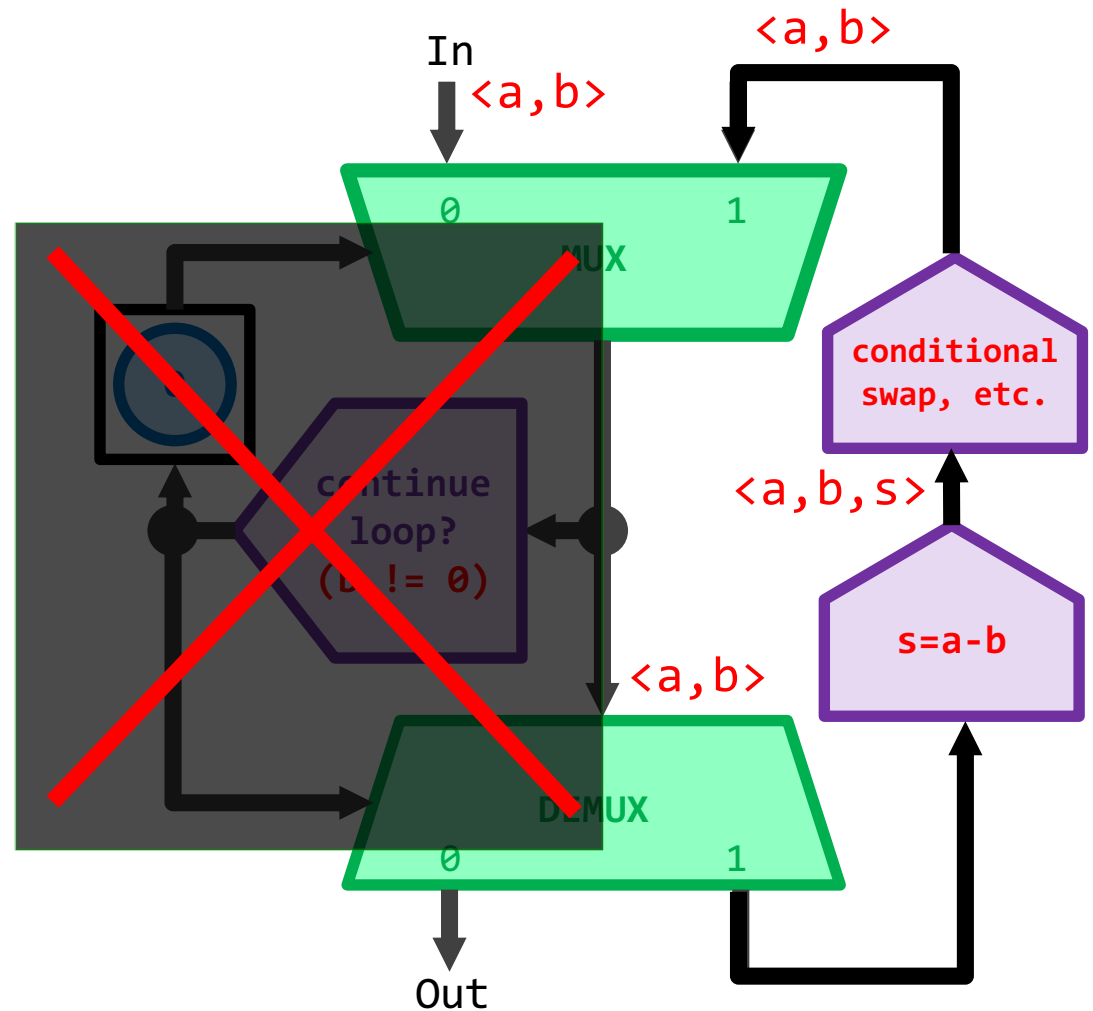
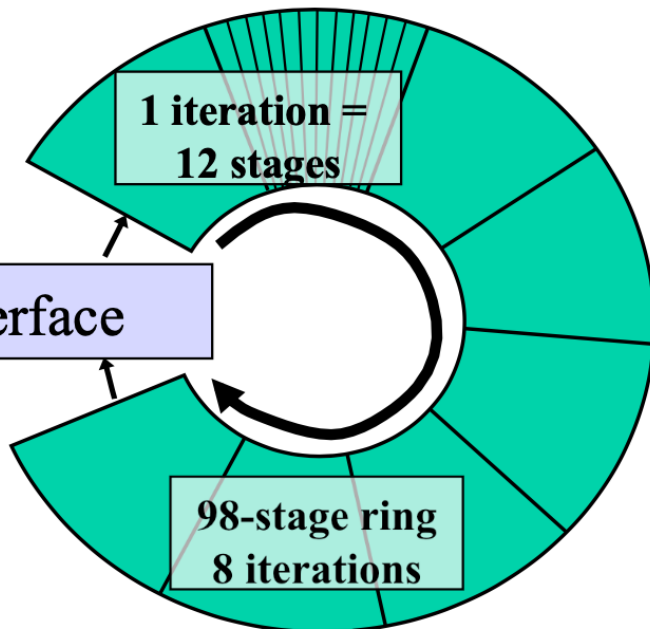
GCD implementation

```
while (b != 0)
```

```

s = a-b
if(s<0)
    swap(a,b)
else
    a = s
    
```

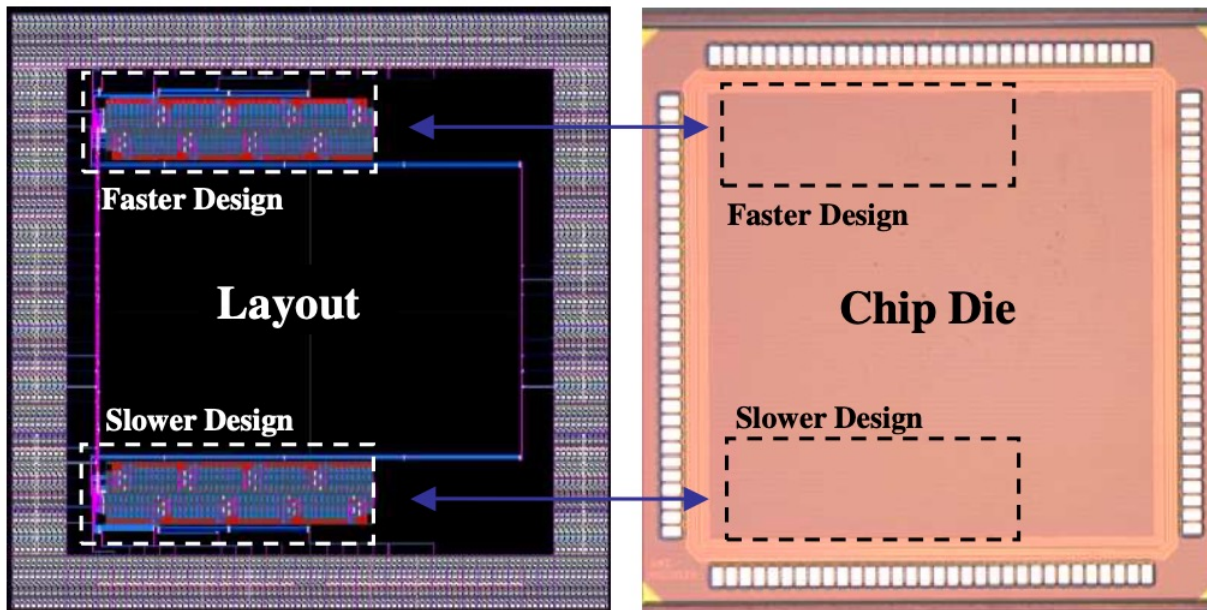
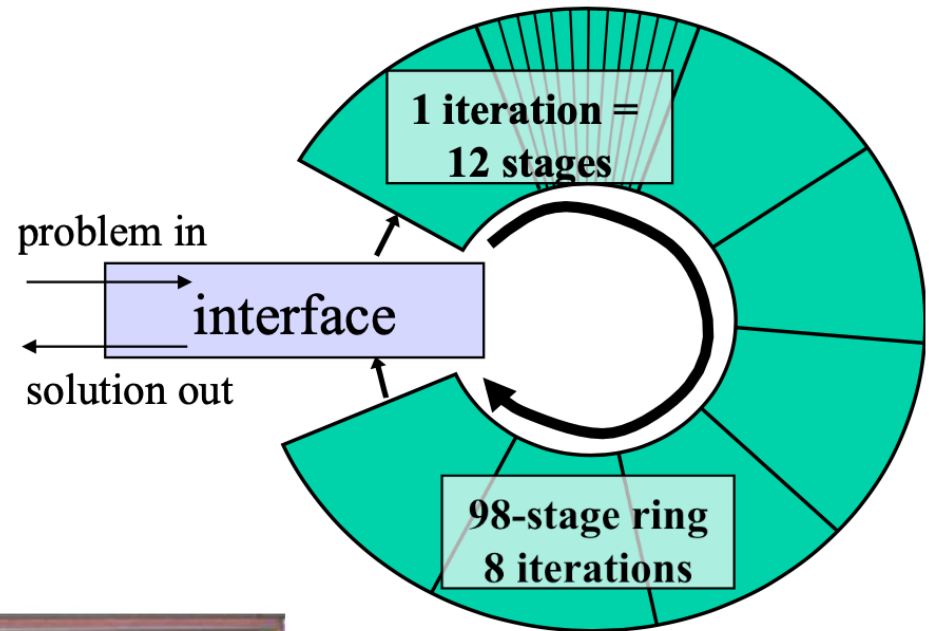
```
return a
```



GCD chip

* Layout and fab:

- 0.13um, standard cell

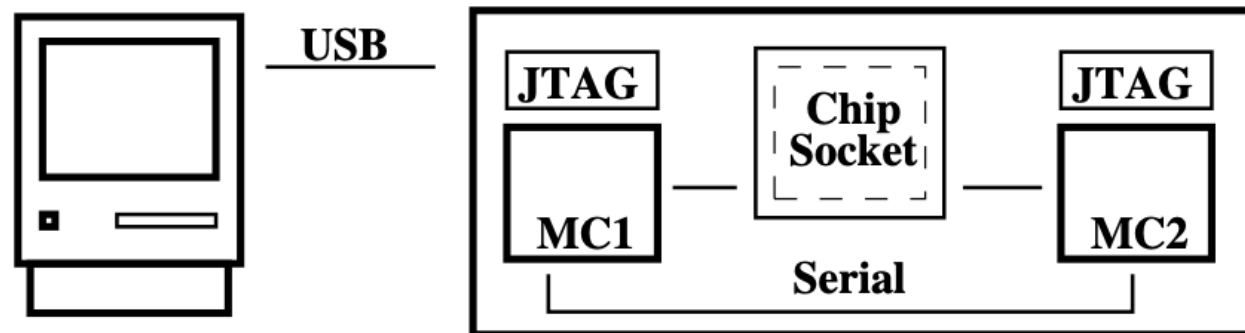


GCD chip

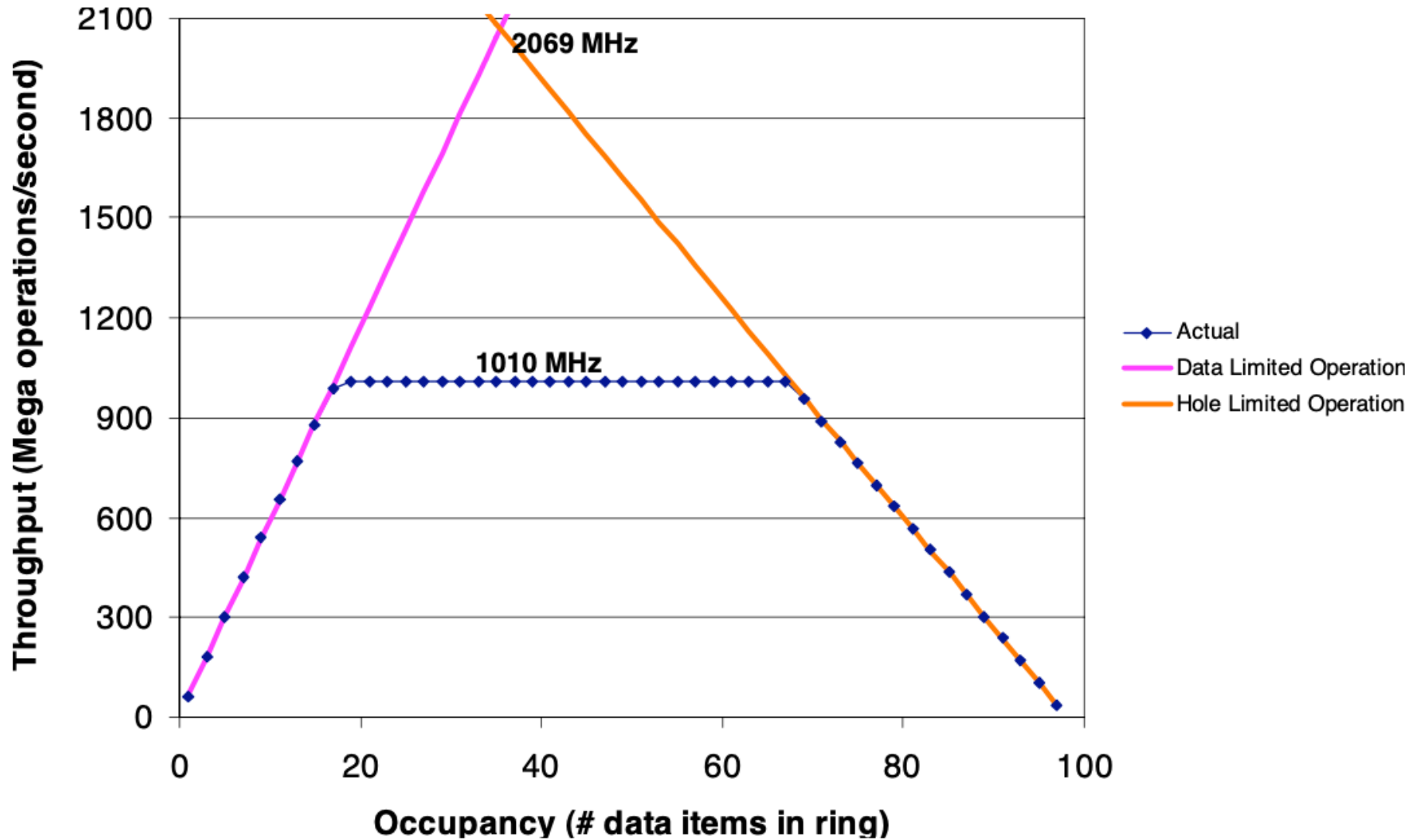
* Testing:

- Full scan for latches
- Combinational test patterns generated (Shi 2005) gave 98% coverage
- Functional tests for timing violations (Gill 2006)

* Evaluation:

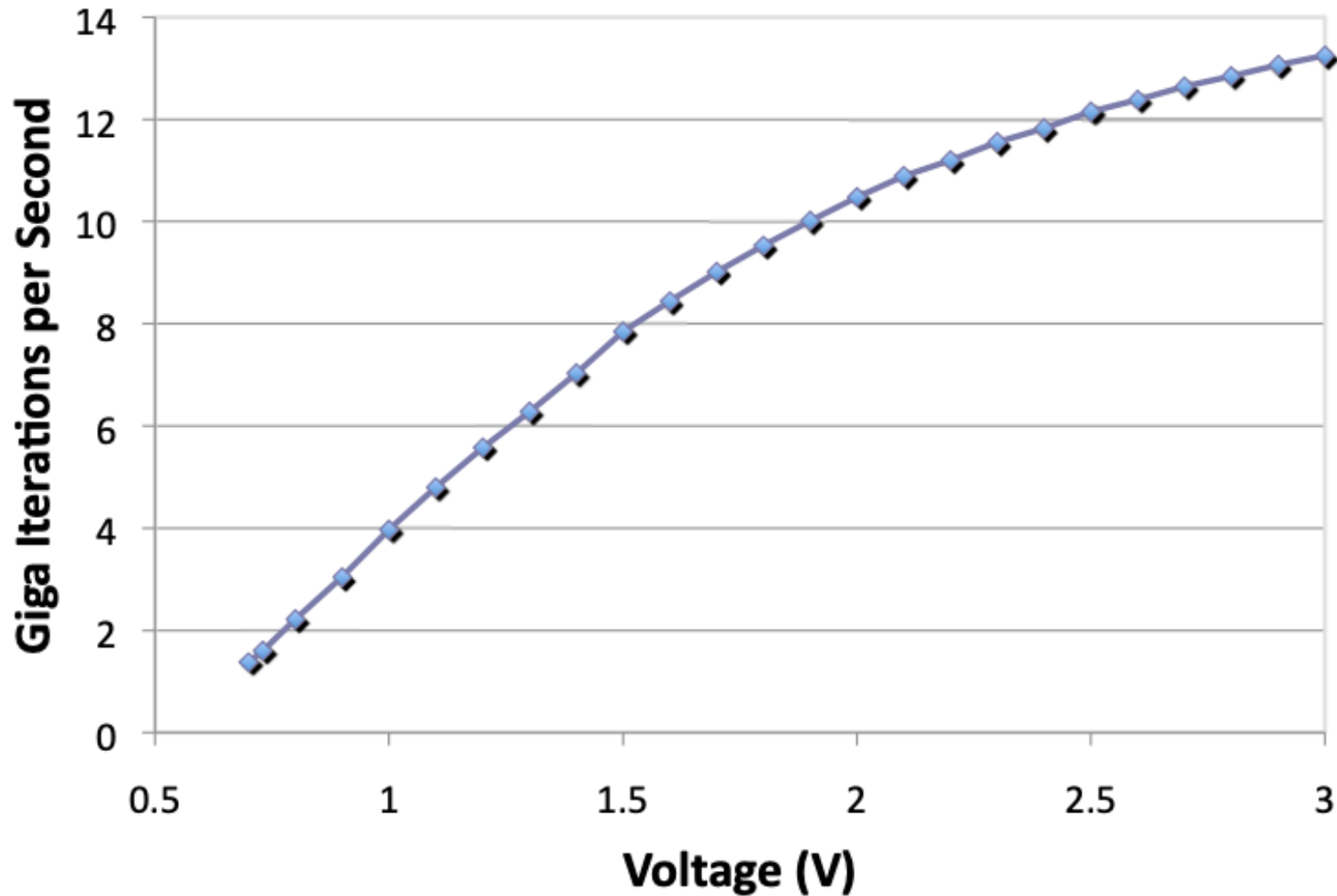


GCD chip: Results



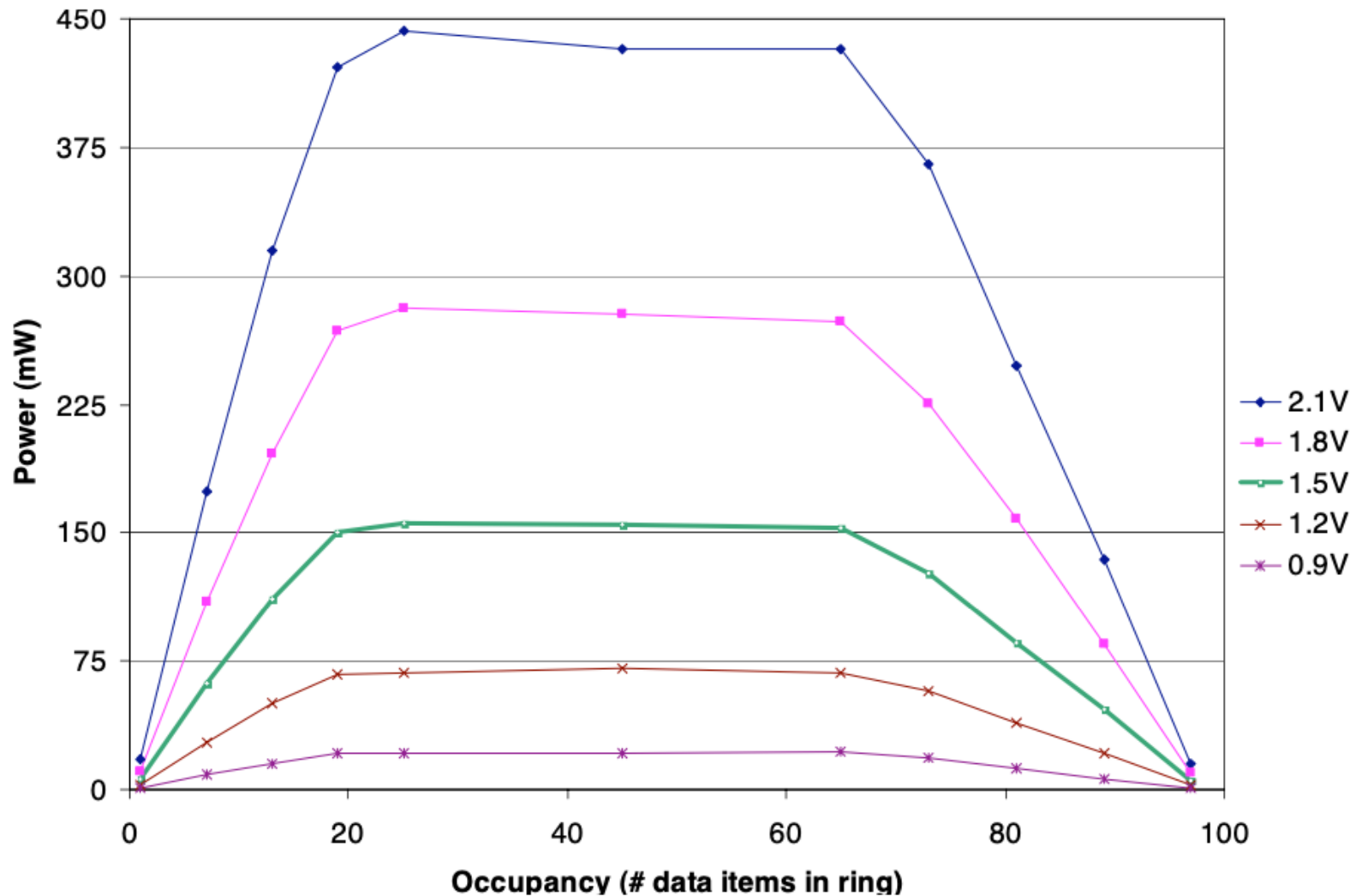
GCD chip: Results

* Operates correctly over a wide voltage range



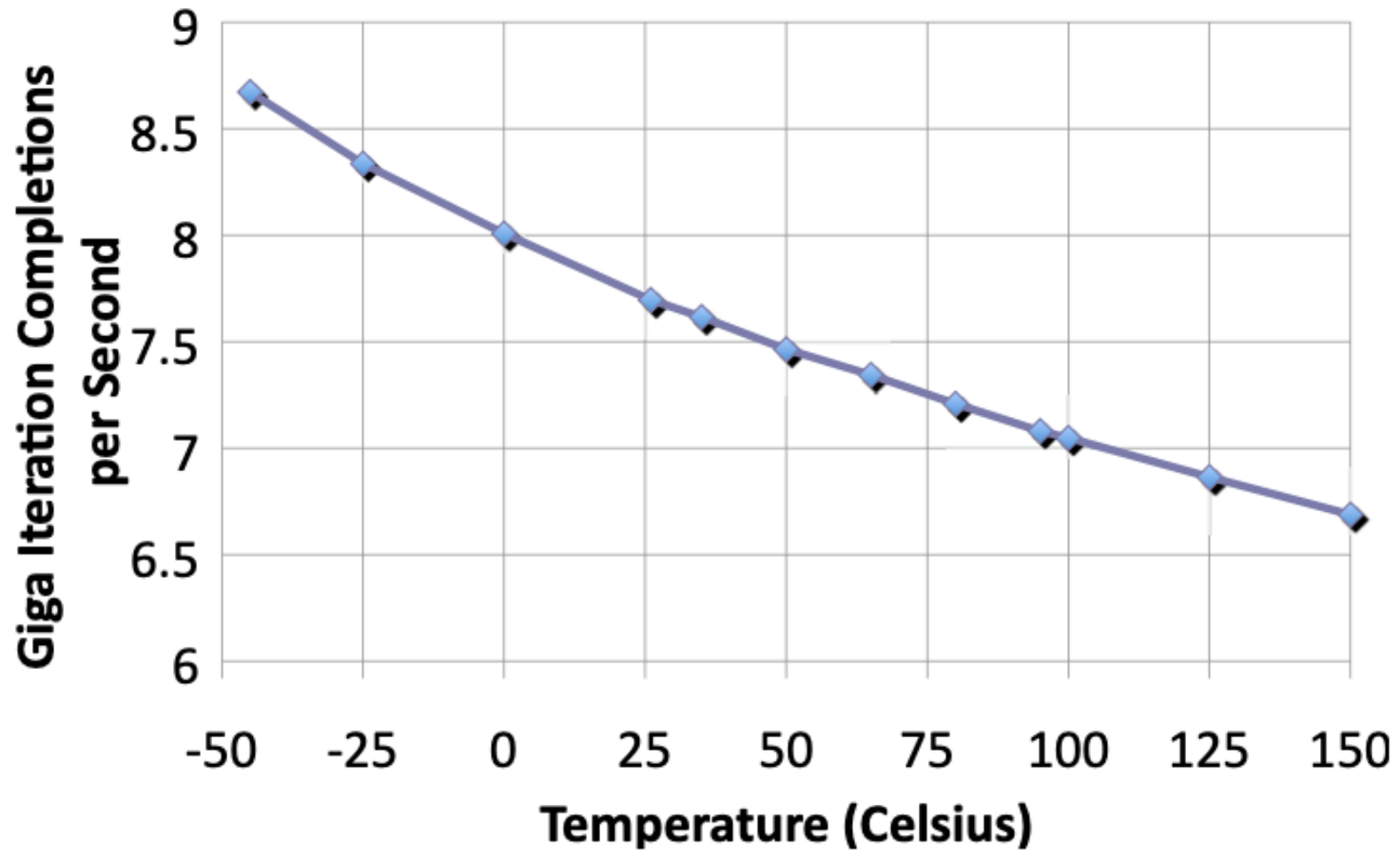
GCD chip: Results

* Impact of voltage variation on power consumption



GCD chip: Results

* Impact of temperature performance



References

- * Feng Shi, Yiorgos Makris, Steven M. Nowick, and Montek Singh. *"Test generation for ultra-high-speed asynchronous pipelines."* ITC 2005.
- * Gennette Gill. *Analysis and Optimization for Pipelined Asynchronous Systems.* PhD thesis. UNC Chapel Hill. 2010.
- * Gennette Gill, A. Agiwal, M. Singh, F. Shi, Y. Makris, *"Low-Overhead Testing of Delay Faults in High-Speed Asynchronous Pipelines."* ASYNC 2006.
- * Montek Singh and Steven Nowick. *"MOUSETRAP: Ultra-High-Speed Transition-Signaling Asynchronous Pipelines."* ICCD 2001.
- * Montek Singh and Steven Nowick. *"MOUSETRAP: High-Speed Transition-Signaling Asynchronous Pipelines."* TVLSI 2007.
- * Gennette Gill, J. Hansen, A. Agiwal, L. Vicci and M, Singh. *"A High-Speed GCD Chip: A Case Study in Asynchronous Design."* ISVLSI 2009.