

# EENG 426/CPSC 459/ENAS 876

## Silicon Compilation

### Single-variable register

Computer Systems Lab

<http://csl.yale.edu/~rajit>

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# Single variable register

Single-variable register:

$$\begin{array}{l} * [\overline{P} \longrightarrow P?x \\ \quad \overline{Q} \longrightarrow Q!x \\ ] \end{array}$$

The process stores one bit of information in local variable  $x$ , which can be read or written.

The environment can communicate on  $P$  or  $Q$  in any order, but mutual exclusion among the two communication actions is guaranteed.

# Handshaking expansion

Handshaking expansion:

$$*[[pti \vee pfi \longrightarrow [pti \longrightarrow x\uparrow \parallel \neg pti \longrightarrow x\downarrow]; po\uparrow; \\ [\neg pti \wedge \neg pfi]; po\downarrow$$
$$[qi \longrightarrow [x \longrightarrow qto\uparrow \parallel \neg x \longrightarrow qfo\uparrow]; [\neg qi]; qto\downarrow, qfo\downarrow \\ ]]$$

Passive protocol on both  $P$  and  $Q$ .

# Handshaking expansion

Write part:

$*[[pti \rightarrow x\uparrow \parallel pfi \rightarrow x\downarrow]; po\uparrow; [\neg pti \wedge \neg pfi]; po\downarrow]$

Production rules:

$pti \mapsto x\uparrow$

$pfi \mapsto x\downarrow$

$pti \wedge x \vee pfi \wedge \neg x \mapsto po\uparrow$

$\neg pti \wedge \neg pfi \mapsto po\downarrow$

# Handshaking expansion

Read part:

$*[[x \wedge qi \longrightarrow qto\uparrow] \neg x \wedge qi \longrightarrow qfo\uparrow]; [\neg qi]; qto\downarrow, qfo\downarrow]$

Production rules:

|                                        |                                 |
|----------------------------------------|---------------------------------|
| $x \wedge qi \mapsto qto\uparrow$      | $\neg qi \mapsto qto\downarrow$ |
| $\neg x \wedge qi \mapsto qfo\uparrow$ | $\neg qi \mapsto qfo\downarrow$ |

# Dual-rail variables

Problem:

$$pti \wedge x \vee pfi \wedge \neg x \mapsto po \uparrow$$

We will need  $x$  and its inverse to turn this into a CMOS implementable circuit.

Solution: introduce the **inverted** version of  $x$  in the HSE

$$\begin{array}{ll} \dots; x \uparrow; \dots & \triangleright \quad \dots; u \downarrow; v \uparrow; \dots \\ \dots; x \downarrow; \dots & \triangleright \quad \dots; v \downarrow; u \uparrow; \dots \end{array}$$

# Dual-rail variables

Write part:

$*[[pti \longrightarrow u\downarrow; v\uparrow] \parallel pfi \longrightarrow v\downarrow; u\uparrow]; po\uparrow; [\neg pti \wedge \neg pfi]; po\downarrow]$

$$pti \mapsto u\downarrow$$

$$pfi \mapsto v\downarrow$$

$$\neg pfi \wedge \neg u \mapsto v\uparrow$$

$$\neg pti \wedge \neg v \mapsto u\uparrow$$

$$pti \wedge v \vee pfi \wedge u \mapsto po\uparrow$$

$$\neg pti \wedge \neg pfi \mapsto po\downarrow$$

# Dual-rail variables

Read part:

$*[[v \wedge qi \longrightarrow qto\uparrow \parallel u \wedge qi \longrightarrow qfo\uparrow]; [\neg qi]; qto\downarrow, qfo\downarrow]$

$$v \wedge qi \mapsto qto\uparrow$$

$$\neg qi \mapsto qto\downarrow$$

$$u \wedge qi \mapsto qfo\uparrow$$

$$\neg qi \mapsto qfo\downarrow$$



# Environment considerations

If the environment can start the next request before the complete handshake is finished:

$$\neg po \wedge v \wedge qi \mapsto qto \uparrow$$

$$\neg po \wedge u \wedge qi \mapsto qfo \uparrow$$

$$\neg qto \wedge \neg qfo \wedge pti \mapsto u \downarrow$$

$$\neg qto \wedge \neg qfo \wedge pfi \mapsto v \downarrow$$

$$\neg qto \wedge \neg qfo \wedge (pti \wedge v \vee pfi \wedge u) \mapsto po \uparrow$$